Representing Dependence

- Represent each *dynamic* instance of a loop as a point in a graph
- Draw arrows from one point to another to represent dependences

 $a[i + 2] = a[i]$ }<br>}

for (i = 0; i < N; i++) {

- Represent each *dynamic* instance of a loop as a point in a graph
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$$
\begin{array}{c}\nfor (i = \ell \\
a[i + 2]\n\end{array}
$$

- Step I: Create nodes, I for each iteration
	- Note: not I for each array location!

$$
\begin{array}{ccc}\n\circ & & \circ \\
\circ & & \circ\n\end{array}
$$

 $\emptyset$ ; i < N; i++) {  $= a[i]$ 



- Represent each *dynamic* instance of a loop as a point in a graph
- Draw arrows from one point to another to represent dependences

• Step 2: Determine which array elements are read and written in each iteration



for (i = 0; i < N; i++) { a[i + 2] = a[i] }

- Represent each *dynamic* instance of a loop as a point in a graph
- Draw arrows from one point to another to represent dependences

• Step 3: Draw arrows to represent dependences



for (i = 0; i < N; i++) { a[i + 2] = a[i] }

# 2-D iteration space graphs

- Can do the same thing for doublynested loops
	- 2 loop counters

for (i = 0; i < N; i++) for (j = 0; j < N; j++) … a[i+1][j] = a[i][j+2] + 1 …



- Can also represent output and anti dependences
	- Use different kinds of arrows for clarity. *E.g.*







• Can we represent dependences in a more compact way?

• Crucial problem: Iteration space graphs are potentially infinite representations!

#### Distance and direction vectors

- Compiler researchers have devised *compressed* representations of dependences
	- Capture the same dependences as an iteration space graph
	- May lose *precision* (show more dependences than the loop actually has)
- Two types
	- **Distance vectors:** captures the "shape" of dependences, but not the particular source and sink
	- **Direction vectors:** captures the "direction" of dependences, but not the particular shape

#### Distance vector

- Represent each dependence arrow in an iteration space graph as a vector
	-



- Distance vector for this iteration space: (2)
	- Each dependence is 2 iterations forward

• Captures the "shape" of the dependence, but loses where the dependence originates

#### 2-D distance vectors

- Distance vector for this graph:
	- $(1, -2)$
	- +1 in the i direction, -2 in the j direction
- Crucial point about distance vectors: they are always "positive"
	- First non-zero entry has to be positive
	- Dependences can't go backwards in time



#### More complex example

#### • Can have multiple distance vectors

$$
\begin{array}{ll}\nfor (i = 0; i < N; i++) \\
for (j = 0; j < N; j++) \\
a[i+2][j] = a[i+1][j+2] + a[i][j]\n\end{array}
$$

 $(2,4)$  $\left( \begin{matrix} 3,4 \end{matrix} \right)$  $(4,4)$  $(1,4)$  $(0,4)$  $\bigodot$  $\bigodot$  $(4,3)$  $(1,3)$  $($ 0,3  $)$  $\bigodot$  $(1,2)$  $\left(3,2\right)$  $(4,2)$  $($ 0,2 $)$  $(2.1)$  $(4,1)$  $(3,1)$  $\left(\left| .\right| \right)$ ( 0,1 )  $(4,0)$  $(2,0)$  $(3,0)$  $($  1,0  $)$  $0,0$ 

#### More complex example

• Can have multiple distance vectors

- Distance vectors
	- $(1, -2)$
	- $(2, 0)$
- Important point: order of vectors depends on order of loops, not use in arrays



$$
\begin{array}{ll}\n\text{for } (i = 0; i < N; i++) \\
\text{for } (j = 0; j < N; j++) \\
a[i+2][j] = a[i+1][j+2] + a[i][j]\n\end{array}
$$

#### Problems with distance vectors

• The preceding examples show how distance vectors can precisely summarize all the dependences in a loop nest using just a small number of distance vectors

- 
- Can't always summarize as easily!
- Running example:





- What are the distance vectors for this code?
	- (1), (2), (3), (4) ...
- Note: we have information about the length of each vector, but not about the source of each vector
	- What happens if we try to reconstruct the iteration space graph?





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#### Direction vectors

- The whole point of distance vectors is that we want to be able to succinctly capture the dependences in a loop nest
	- But in the previous example, not only did we add a lot of extra information, we still had an infinite number of distance vectors
- Idea: summarize distance vectors, and save only the *direction* the dependence was in
	- $(2, -1) \rightarrow (+, -)$
	- $(0, 1) \rightarrow (0, +)$
	- $(0, -2) \rightarrow (0, -)$  (can't happen; dependences have to be positive)
	- Notation: sometimes use '<' and '>' instead of '+' and '-'

Why use direction vectors?

- Direction vectors lose a lot of information, but do capture some useful information
	- Whether there is a dependence (anything other than a '0' means there is a dependence)
	- Which dimension and direction the dependence is in
- Many times, the only information we need to determine if an optimization is legal is captured by direction vectors
	- Loop parallelization
	- Loop interchange