

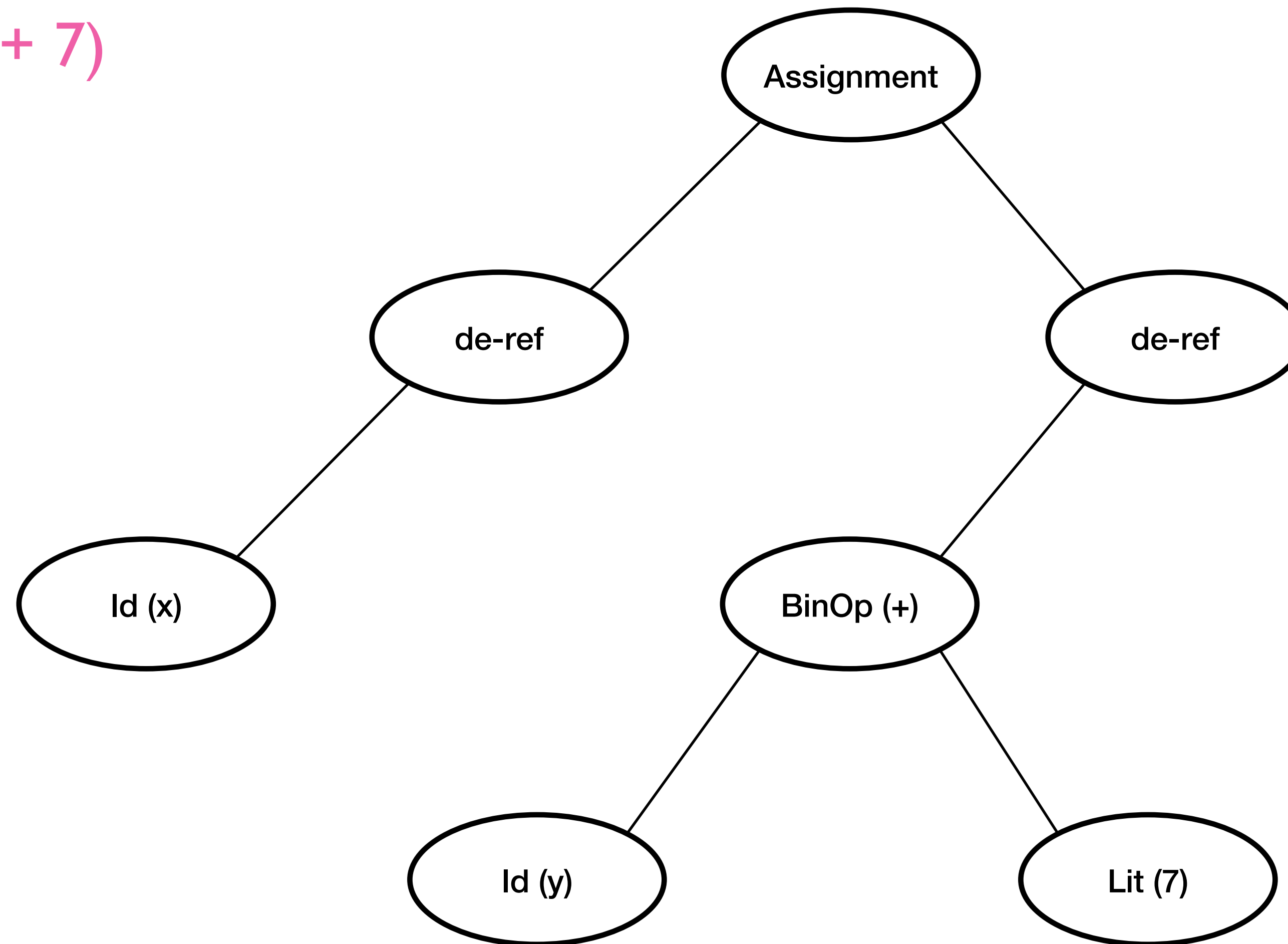
Pointer Codegen Example

code generation (assembly)

- Code generation in assembly is easy: keep the same CodeObject, but switch whether temporary is an l-val or an r-val

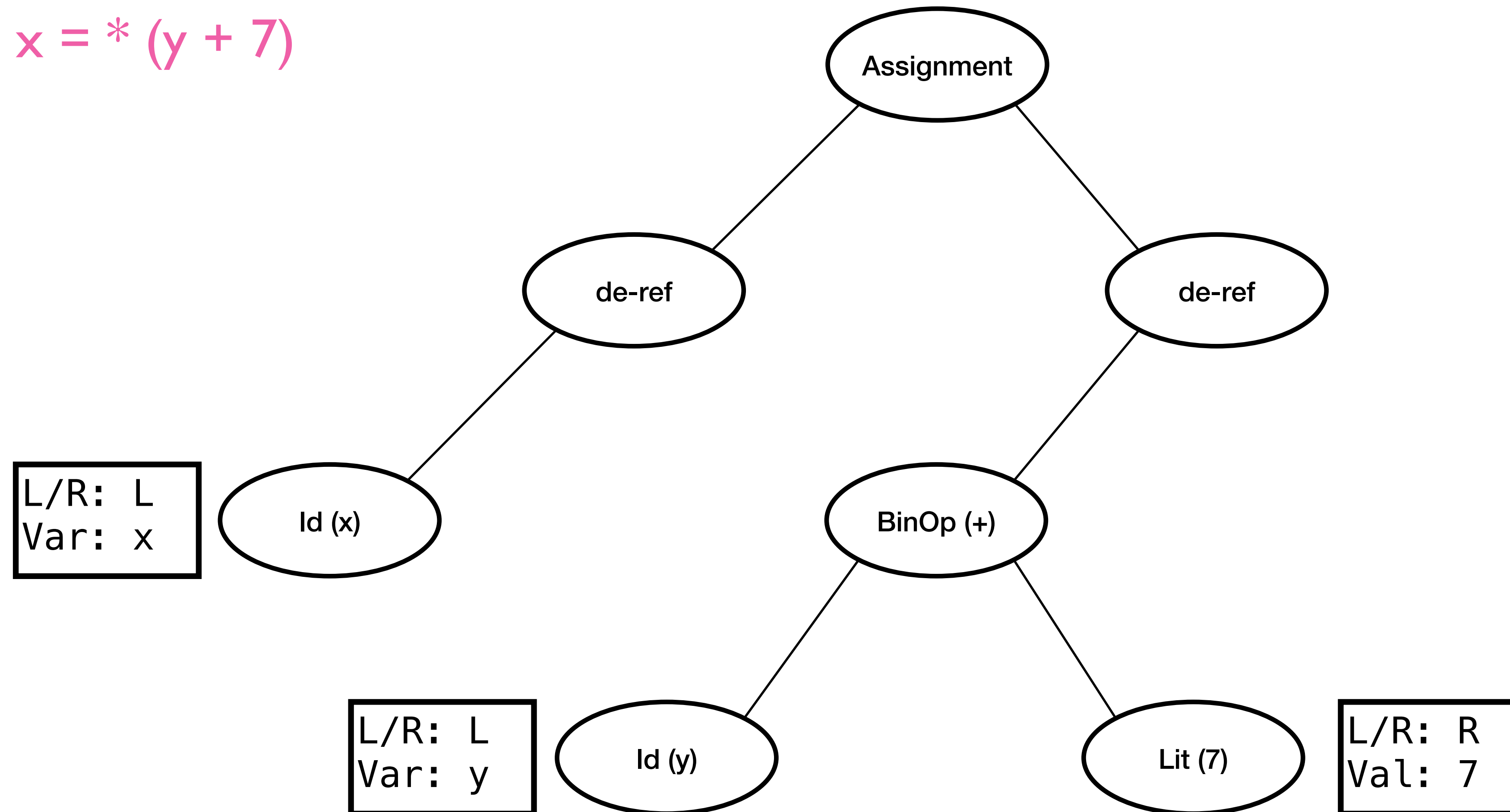
code generation (assembly)

- $*x = *(y + 7)$



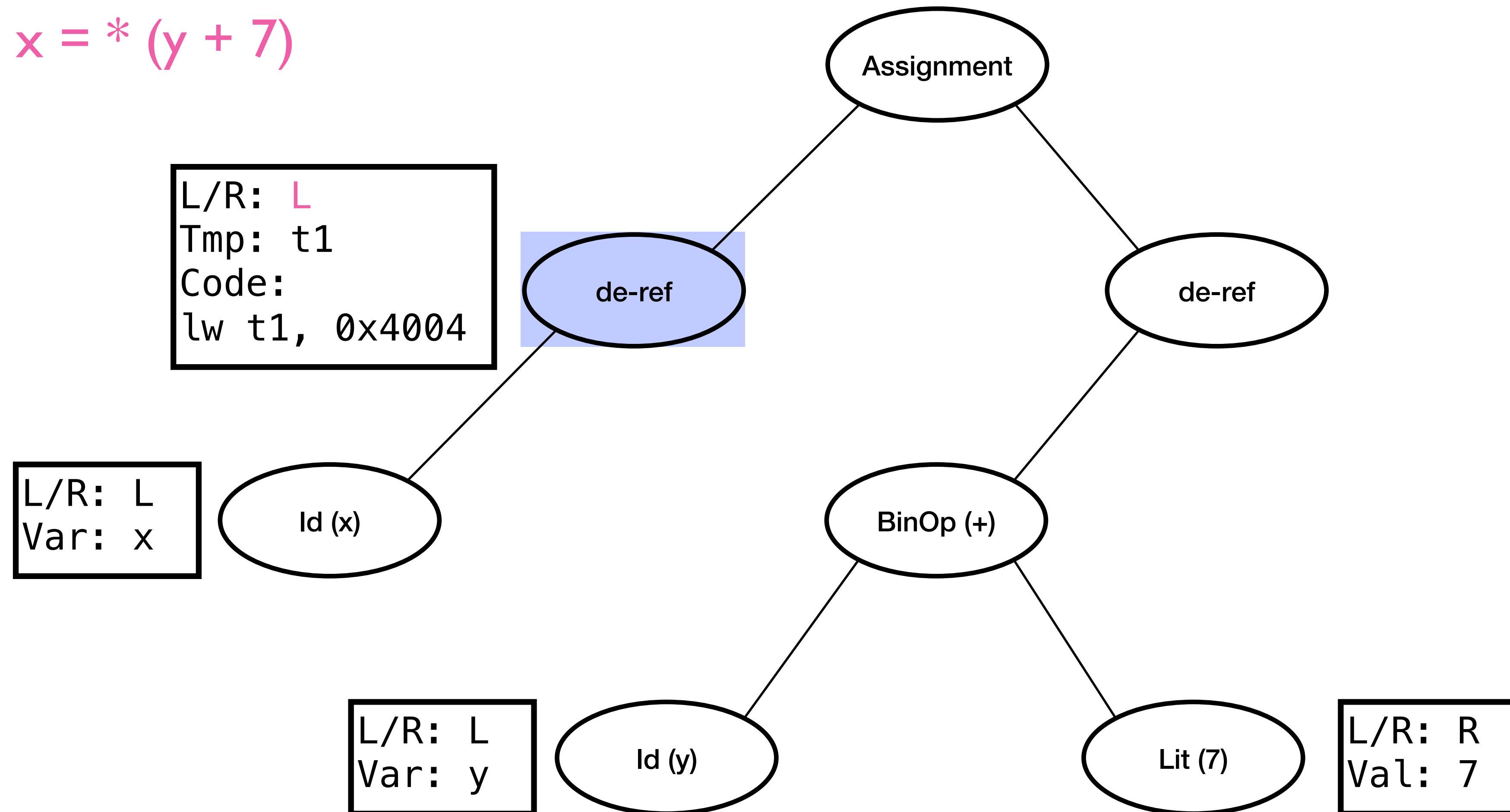
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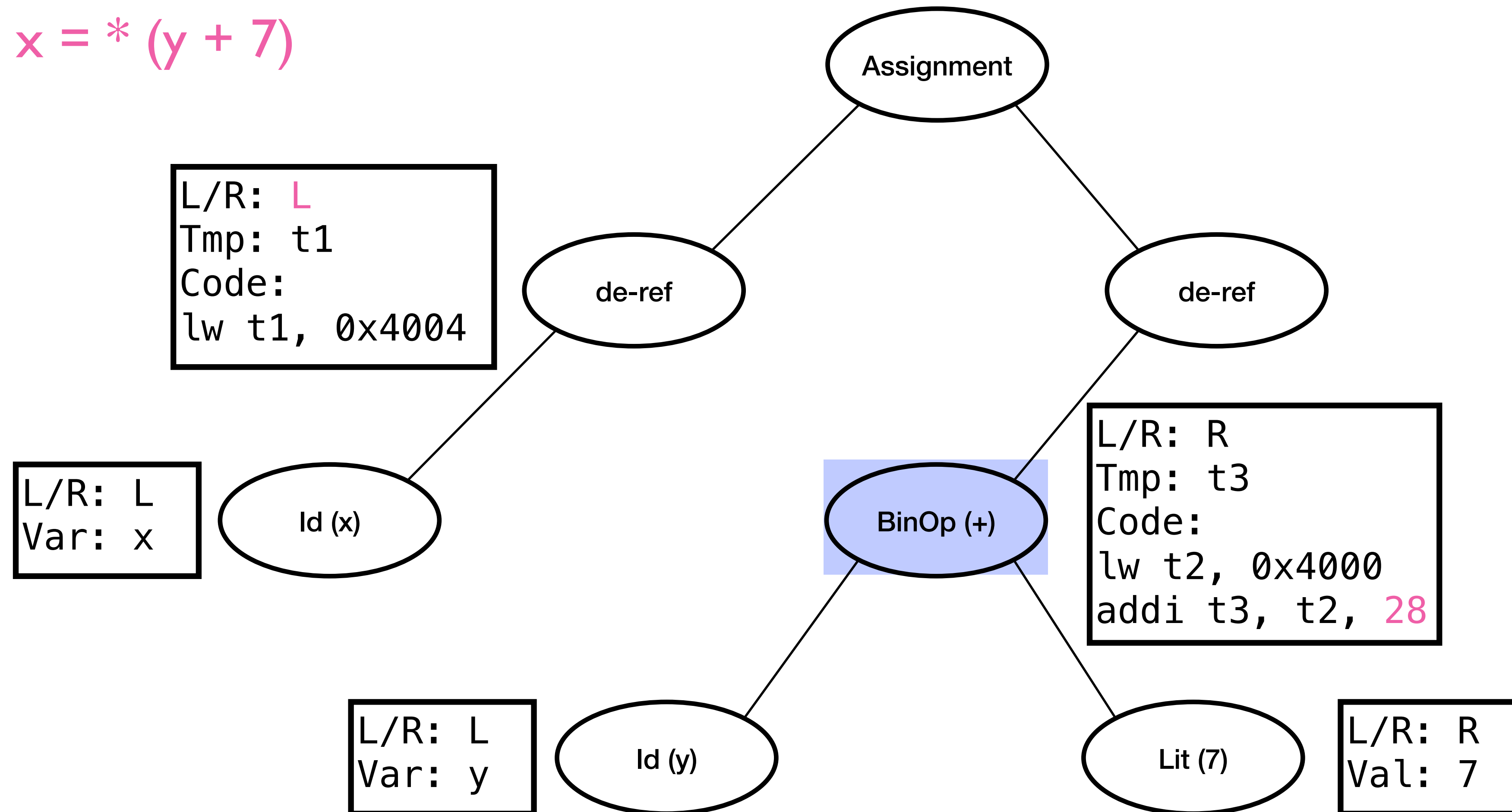
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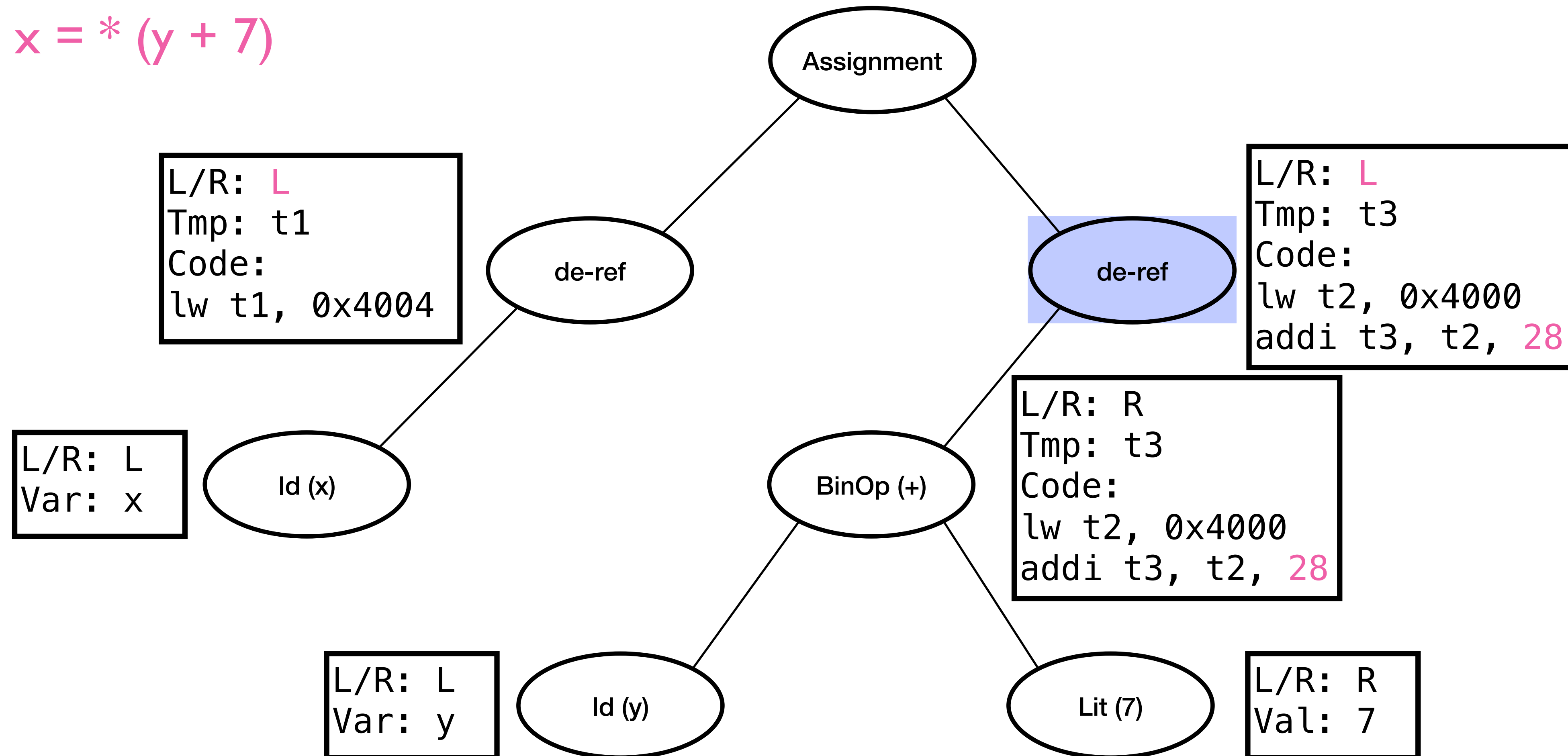
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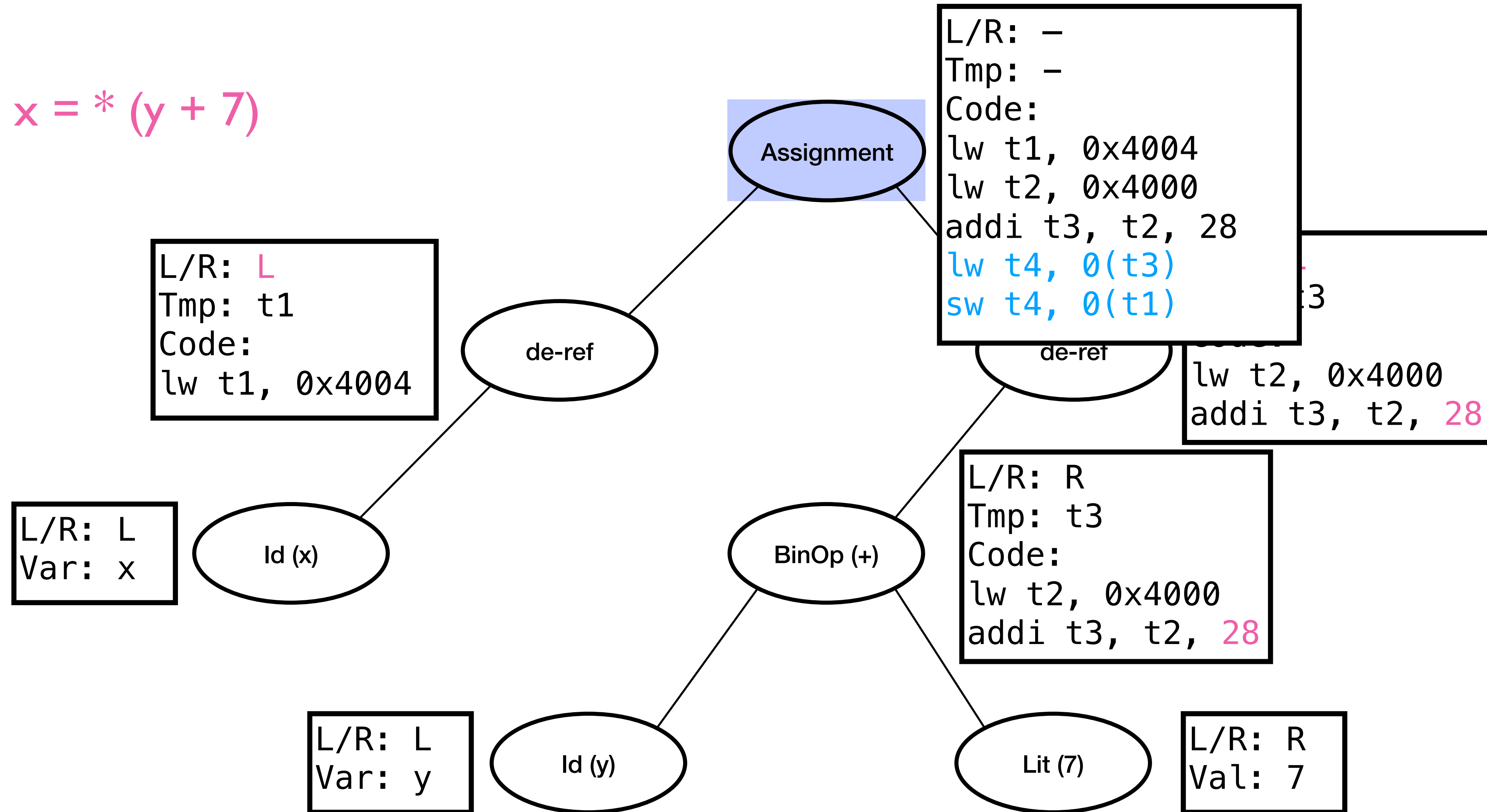
code generation (assembly)

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code generation (assembly)

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code generation (IR)

- Code generation for IR is similar
- Track whether IR temporary holds an l-value or an r-value (e.g., use '\$' as prefix for r-value, '@' as prefix for l-value)
- Introduce two new IR nodes:
 - **ADDROF** *a, b* : store the address of operand *b* in *a* (if *b* is a variable, *a* holds the variable address; if *b* is an l-value temporary, *a* is the temporary, just as an r-value)
 - **DEREF** *a, b* : store the value of operand *b* in *a* as an address (if *b* is a variable or an l-value temporary, load from *b* and store the result in *a* as an l-value; if *b* is an r-value temporary, *a* is the temporary, just as an l-value)

register allocation

- Now that we have pointers, we have aliasing!
- Simple solution: treat all locals/globals as aliased to each other: cannot stay in registers. Write back on every store, free after every load
- Slightly more complicated: only variables that have ever had an ADDROF operation applied to them can be aliased
- More complex: perform **pointer analysis** (stay tuned!)

next: memory allocation